OPTIMIZED IMPLEMENTATION OF FFT PROCESSOR FOR OFDM SYSTEMS

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ABSTRACT
Due to the advanced VLSI technology, Fast Fourier Transform (FFT) and Inverse Fast Fourier Transform (IFFT) has been applied to wide field of digital signal processing (DSP) applications. For modern communication systems, the FFT/IFFT is very important in the OFDM. Because of FFT/IFFT is the key computational block to execute the baseband multicarrier demodulation and modulation in the OFDM system. The Fast Fourier Transform (FFT) and its inverse transform (IFFT) processor are key components in OFDM systems. An optimized implementation of the 8-point FFT processor with radix-2 algorithm in R2MDC architecture is the processing element. This butterfly- Processing Element (PE) used in the 8-FFT processor reduces the multiplicative complexity by using a real constant multiplication in one method and eliminates the multiplicative complexity by using add and shift operations in this proposed method.

KEYWORDS: Cooley-Tukey, R2MDC, FFT, VHDL, OFDM.

I. INTRODUCTION
In the last few years, wireless communication plays a significant role in people’s life. The development of wireless communication science and technology was extremely fast due to VLSI technology. Many complicated design of communication system become feasible. There is a rapidly growing demand for high speed, efficient and reliable digital communication system with the purpose of support multimedia transmission such as high-quality voice, video, etc. Orthogonal Frequency Division Multiplexing (OFDM) technology is an effective modulation scheme to meet the demand. This is an advanced modulation technique that effectively expands channel utilization and reduces inter-symbol interference (ISI) and inter-carrier interference (ICI) caused by multi-path effect. Multiple-input multiple-output (MIMO) signal processing technique has been utilized in combination with OFDM for next generation wireless communication system to enhance link throughput as well as the robustness of transmission over frequency selective fading channel. FFT algorithms are based on the fundamental principle of decomposing the computation of the DFT of a sequence of length N into successively smaller DFT using the symmetry, periodicity, compressibility and expansibility properties of WN. The manner in which this principle is implemented leads to a variety of different algorithms, all with comparable improvements in computational speed. It can reduce the computational complexity from O (N2) to O (Nlog2N), and the regularity of the algorithm makes it suitable for VLSI implementation. Due to its advantage, it can be efficiently employed in the modern Orthogonal Frequency Division Multiplexing (OFDM) system. The OFDM signal is made up of many orthogonal characters, and each individual carrier is digitally modulated with a relatively slow symbol rate. This method has distinct advantages in multipath propagation because, in comparison with the single carrier method at the same transmission rate, more time is needed to transmit a symbol. The 256QAM modulation modes are used, and the modulation is adapted to the specific transmission requirements. Whereas 64-QAM modulation mode is used for Wi-Fi. Transmission rates of up to 75 Mbit/s are possible at 1GB speed. The 802.16e is a standard of Wimax in the frequency range of up to 6 GHz with the objective of allowing mobile applications and even roaming. In addition, the number
of carriers can vary over a wide range depending on permutation zone and FFT base (128, 512, 1024, and 2048). Based on our development, not only a dedicated FFT/IFFT module can be easily prototyped for fast system verification, but also the resulting compiler can be used as a basis for more advanced research in the future.

This work can be organized as: section (I) explains the concept of OFDM, FFT processor necessity. Section (II) FFT algorithm with the help of radix point. Section (III) architecture of r2mdc and add shift method section (IV) simulation results section(V) synthesis report section (VI) conclusion

II. CONVENTIONAL FFT ALGORITHM

Discrete Fourier Transform (DFT) is widely applied to the analysis, design, and the implementation of discrete-time signal processing algorithms and communications. The computational complexity of direct evaluation is rather high so that it is hard to meet high performance and low cost design goal. Therefore, a fast DFT algorithm is required. Since the early paper proposed by Cooley and Tukey in 1965 [3], they provided a lot of ways to reduce the computational complexity from $O(N^2)$ to $O(N \log_2 N)$. After that, the algorithms that can reduce the computational complexity are generally called fast Fourier transform (FFT) algorithms.

The fast Fourier transform (FFT) algorithm, together with its many successful applications, represents one of the most important advancements in scientific and engineering computing in this century. The wide usage of computers has been instrumental in driving the study of the FFT, and a very large number of articles have been written about the algorithm over the past thirty years. FFT algorithms are based on the fundamental principle of decomposing the computation of the discrete Fourier transform of a sequence of length N into successively smaller discrete Fourier transforms. The regularity of the algorithms makes them suitable for VLSI implementation.

Decomposing is an important role in the FFT algorithms. There are two decomposed types of the FFT algorithm. One is decimation-in-time (DIT), and the other is decimation-in-frequency (DIF) shown in fig 1. The difference between these two types is in the input 6 and output data ordering in signal flow graph (SFG). The DIT algorithm means that the time sequence is decomposed into small subsequence, and the DIF algorithm decomposes the frequency sequence. In addition, there is no difference in computational complexity between these two types of FFT algorithm. Since the low computational complexity of FFT algorithms is desired for high speed consideration in VLSI implementation, here we discuss the computational complexity of different algorithms.
Basic Radix-2 butterfly processor shown in Fig. 2. consists of a complex adder and complex subtraction. Besides that, an additional complex multiplier for the twiddle factors WN is implemented. The complex multiplication with the twiddle factor requires four real multiplications and two add/subtract operations. Three twiddle factor values are used i.e. c, c+s, c-s.

III. PROPOSED ARCHITECTURE

A. R2MDC ARCHITECTURE

One of the most straightforward approaches for pipeline implementation of radix-2 FFT algorithm is Radix-2 Multi-path Delay Commutator (R2MDC) architecture. It’s the simplest way to rearrange data for the FFT/IFFT algorithm. The input data sequence are broken into two parallel data stream flowing forward, with correct distance between data elements entering the butterfly scheduled by proper delays. 8-point FFT in R2MDC architecture is shown in Fig. 4.

At each stage of this architecture half of the data flow is delayed via the memory (Reg) and processed with the second half data stream. The delay for each stage is 4, 2, and 1 respectively. The total number of delay elements is 4 + 2 + 2 + 1 + 1 = 10. In this R2MDC architecture, both Butterflies (BF) and multipliers are idle half the time waiting for the new inputs. The 8-point FFT/IFFT processor has one multiplier, 3 of radix-2 butterflies, 10 registers (R) (delay elements) and 2 switches (S).

\[ x_{3}x_{2}x_{1}x_{0} \]
\[ x_{7}x_{6}x_{5}x_{4} \]

**Fig.4. R2MDC ARCHITECTURE**

B. ADD AND SHIFT METHOD

Another method proposed eliminates the non-trivial complex multiplication with the twiddle factors (W8 1, W8 3) and implements the processor without complex multiplication. The proposed butterfly processor performs the multiplication with the trivial factor W8 2=-j by switching from real to imaginary part and imaginary to real part, with the factor W8 0 by a simple cable. With the non-trivial factors W1, W3 the processor realize the multiplication by the factor 1/√2 using hard wired shift-and-add operation as shown in Fig. 5.

\[ x(3)x(2)x(1)x(0) \]
\[ x(7)x(6)x(5)x(4) \]

**Fig.5. Butterfly processor with no complex multiplication**
IV. ANALYSIS AND DESIGN OF FFT USING VHDL

A. Simulation methodology

VHDL is a language for describing digital electronic systems. It arose out of the United States government’s very high speed integrated circuits (VHSIC) program, in the course of this program; it became clear that there was a need for a standard language for describing the structure and function of integrated circuits (IC’s). Hence the VHSIC hardware description language (VHDL) was developed. Modeling for simulation and synthesis is a vital part of a range of levels of abstraction, from gate levels up to algorithmic and architectural levels. It will continue to play an important role in the design of silicon-based systems for some time to come. Very high speed integrated circuit hardware description language (VHDL) can be used to model digital systems and introduce some of the basic concepts underlying the language.

The ModelSim-Altera simulator compiles the testbench and the netlist (multiplier.vho), annotates the SDF data (in multiplier_vhd.sdo), and runs the simulation for the specified time. A waveform window within the ModelSim-Altera simulator is invoked that shows the expected and actual results of the multiplier. The expected and actual results are also checked in the test bench, and messages that show whether or not the results match are displayed in the simulator’s console window. The data output of the multiplier module changes with a delay after the clock edge because the SDF data is annotated in the gate-level timing simulation.

B. Simulation result

The simulation results are given below.

![Simulation result of R2MDC architecture](image1.png)

This simulation shows the radix-2 operation. Two inputs are selected with the help of demultiplexer which is present inside the buffer block. Addition and subtraction of two inputs with a twiddle factor multiplication. It produces final fft output.

The VHDL Program is compiled in Model sim v6.4a to generate the butterfly processor operation with the help of R2MDC architecture, Add and shift method to know the design functionality and for power, area analysis Xilinx is used.
Fig. 7. Simulation result of add and shift method

This simulation shows that the input after addition and subtraction of radix-2 process multiplied with a real number of complex variable. It produces the same FFT output.

V. POWER AND AREA ANALYSIS OF R2MDC AND ADD SHIFT METHOD

Table 1 Power Analysis

<table>
<thead>
<tr>
<th>Architecture</th>
<th>Synthesis Power</th>
<th>Multiplication</th>
</tr>
</thead>
<tbody>
<tr>
<td>R2MDC Architecture</td>
<td>0.323</td>
<td>3</td>
</tr>
<tr>
<td>Add Shift Method</td>
<td>1.517</td>
<td>0</td>
</tr>
</tbody>
</table>

Table 2 Area Analysis

<table>
<thead>
<tr>
<th>Logic Utilization</th>
<th>Used</th>
<th>Available</th>
<th>Utilization</th>
</tr>
</thead>
<tbody>
<tr>
<td>Number of 4 input LUTs</td>
<td>148</td>
<td>3,840</td>
<td>3%</td>
</tr>
</tbody>
</table>

<table>
<thead>
<tr>
<th>Logic Distribution</th>
<th>Used</th>
<th>Available</th>
<th>Utilization</th>
</tr>
</thead>
<tbody>
<tr>
<td>Number of occupied Slices</td>
<td>80</td>
<td>1,920</td>
<td>4%</td>
</tr>
<tr>
<td>Number of Slices containing only related logic</td>
<td>80</td>
<td>80</td>
<td>100%</td>
</tr>
<tr>
<td>Number of Slices containing unrelated logic</td>
<td>0</td>
<td>80</td>
<td>0%</td>
</tr>
<tr>
<td>Total Number of 4 input LUTs</td>
<td>157</td>
<td>3,840</td>
<td>4%</td>
</tr>
<tr>
<td>Number used as logic</td>
<td>148</td>
<td></td>
<td></td>
</tr>
<tr>
<td>Number used as a route-thru</td>
<td>9</td>
<td></td>
<td></td>
</tr>
<tr>
<td>Number of bonded IOBs</td>
<td>34</td>
<td>141</td>
<td>24%</td>
</tr>
</tbody>
</table>

With the simulation output a vcd file is created in the transcript of modelsim and it is synthesis in Xilinx to show the power and area analysis for r2mdc and add shift method.

<table>
<thead>
<tr>
<th>Logic Utilization</th>
<th>Used</th>
<th>Available</th>
<th>Utilization</th>
</tr>
</thead>
<tbody>
<tr>
<td>Number of Slice Latches</td>
<td>16</td>
<td>3,840</td>
<td>1%</td>
</tr>
<tr>
<td>Number of 4 input LUTs</td>
<td>85</td>
<td>3,840</td>
<td>2%</td>
</tr>
</tbody>
</table>

<table>
<thead>
<tr>
<th>Logic Distribution</th>
<th>Used</th>
<th>Available</th>
<th>Utilization</th>
</tr>
</thead>
<tbody>
<tr>
<td>Number of occupied Slices</td>
<td>52</td>
<td>1,920</td>
<td>2%</td>
</tr>
<tr>
<td>Number of Slices containing only related logic</td>
<td>52</td>
<td>52</td>
<td>100%</td>
</tr>
<tr>
<td>Number of Slices containing unrelated logic</td>
<td>0</td>
<td>52</td>
<td>0%</td>
</tr>
<tr>
<td>Total Number of 4 input LUTs</td>
<td>87</td>
<td>3,840</td>
<td>2%</td>
</tr>
</tbody>
</table>
VI. CONCLUSION

In this paper, two pipeline-based FFT Architectures are proposed. Our proposed variable-length FFT processor that are suitable for various MIMO OFDM-based communication systems, such as IEEE 802.11n, and IEEE 802.16 WiMAX, can perform 256/128/64/16/8-point with 1-4 data sequences. The proposed 8-point FFT processor is used for IEEE 802.11n, and the proposed 8-point FFT processor is used for IEEE 802.11n and IEEE 802.16. To reduce computational complexity and increase hardware utility, we adopt different radix FFT algorithms and multiple-path delay commutates FFT architecture in our processors. The multiple-path delay commutator FFT architectures require fewer delay elements and different radix FFT algorithms require fewer complex multiplications. The processor Implementation are fabricated using UMC 0.18 $\mu m$ process and their area are 1.5162 2 mm and 2.5122 2 mm.

REFERENCES

[5]. IEEE 802.15 WPAN Task Group 3c (TG3c) Millimeter Wave Alternative PHY Friday, 5 February 2010,
[6]. Hardware Implementation Low Power High Speed FFT Core The International Arab Journal of Information Technology, Vol. 6, No. 1, January 2009

Appendix
The program of R2mdc architecture

```vhdl
library ieee;
use ieee.std_logic_1164.all;
use ieee.std_logic_arith.all;
use ieee.std_logic_signed.all;
use ieee.std_logic_math.all;

--signal x:std_logic;
begin
  process(a,r)
  begin
    c<=a+r;
  end process;
end entity;
```

adder

library ieee;
use ieee.std_logic_1164.all;
use ieee.std_logic_arith.all;
use ieee.std_logic_signed.all;
--use ieee.std_logic_math.all;

entity adder is
  port(a,r,in std_logic_vector(3 downto 0));
c:out std_logic_vector(3 downto 0));
end entity;
architecture adder of adder is
begin
  process(a,r)
  begin
    c<=a+r;
  end process;
end process;
```
library ieee;
use ieee.std_logic_1164.all;
use ieee.std_logic_arith.all;
use ieee.std_logic_unsigned.all;

entity butterfly is
    port(x1:in std_logic_vector(3 downto 0);
        x2:in std_logic_vector(3 downto 0);
        c1:in std_logic;
        y1,y2:out std_logic_vector(3 downto 0));
end entity;

architecture butterfly of butterfly is

component adder is
    port(a,r:in std_logic_vector(3 downto 0);
        c:out std_logic_vector(3 downto 0));
end component;

component subtractor is
    port(a,s:in std_logic_vector(3 downto 0);
        c:out std_logic_vector(3 downto 0));
end component;

component mux is
    port(sel:in std_logic;
        a:in std_logic_vector(3 downto 0);
        b:in std_logic_vector(3 downto 0);
        y:out std_logic_vector(3 downto 0));
end component;

signal s1,s2,s3,s4:std_logic_vector(3 downto 0);
signal o1,o2:std_logic_vector(3 downto 0);
signal z,q:std_logic_vector(3 downto 0);
begin
    process(x1,x2,c1)
    begin
        o1<=o1*00000010;
        o2<=o2*00000010;
        y1<=o1;
        y2<=o2;
    end process p;
end architecture;

DEMUX

library ieee;
use ieee.std_logic_1164.all;

entity demux is
    port(i:in std_logic_vector(3 downto 0);
        sel2:in std_logic_vector(1 downto 0);
        y0:out std_logic_vector(3 downto 0);
        y1:out std_logic_vector(3 downto 0));
end demux;

architecture Behavioral of demux is
begin
    process(i,sel2)
    begin
        case sel2 is
            when "00" => y0 <= i;
            when "11" => y1 <= i;
            when others => --out0 <= bitin;
        end case;
    end process;
end architecture;
INTERBLOCK

library ieee;
use ieee.std_logic_1164.all;
use ieee.std_logic_arith.all;
use ieee.std_logic_unsigned.all;
--use ieee.std_logic_math.all;

dentity interblock is
  port(j,k: in std_logic_vector(3 downto 0);
    sel:in std_logic;
    y0:out std_logic_vector(3 downto 0);
    y1:out std_logic_vector(3 downto 0));
end interblock;
architecture fft2 of interblock is
begin
  process(j,k)
    variable s:std_logic_vector(3 downto 0);
  begin
    s:="0001";
    case sel is
    when '0' =>
      y0<=j+s;
    when '1' =>
      y1<=k;
    when others =>
      end case;
  end process;
end architecture;

MULTIPLIER

library ieee;
use ieee.std_logic_1164.all;
use ieee.std_logic_arith.all;
use ieee.std_logic_unsigned.all;
--use ieee.std_logic_math.all;

dentity multi is
  port(inp,w: in std_logic_vector(3 downto 0);
    oup:out std_logic_vector(3 downto 0));
end multi;
architecture fft3 of multi is
signal sin:std_logic_vector(7 downto 0);
signal si:std_logic_vector(7 downto 0);
--signal k:std_logic_vector(7 downto 0);
begin
  oup<=sin(3 downto 0);
end fft3;

MULTIPLEXER

library ieee;
use ieee.std_logic_1164.all;
use ieee.std_logic_arith.all;
use ieee.std_logic_unsigned.all;

dentity mux is
  port(sel:in std_logic;
    a:in std_logic_vector(3 downto 0);
    b:in std_logic_vector(3 downto 0);
    y:out std_logic_vector(3 downto 0));
end entity;
architecture mux of mux is
begin
  process(a,b,sel)
    begin
      case sel is
      when '0' =>
        y<=a;
      when others =>
        y<=b;
      end case;
  end process;
end mux;

SUBTRACTOR

library ieee;
use ieee.std_logic_1164.all;
use ieee.std_logic_arith.all;
use ieee.std_logic_signed.all;

entity subtractor is
  port(a,s:in std_logic_vector(3 downto 0);
    c:out std_logic_vector(3 downto 0));
end entity;
architecture subtractor of subtractor
begin
  process(a,s)
  begin
    c<=a-s;
  end process;
end architecture;

TOP MODULE

library ieee;
use ieee.std_logic_1164.all;
use ieee.std_logic_arith.all;
entity top_module is
begin
  process(a,s)
  begin
    c<=a-s;
  end process;
end architecture;
The program for add and shift method

library ieee;
use ieee.std_logic_1164.all;
use ieee.std_logic_unsigned.all;
entity append is
  port(a:in std_logic_vector(21 downto 0);
       c:out std_logic_vector(22 downto 0));
end entity;
architecture beh of append is
begin
  process(a)
  begin
    c<=(ord('0')&a);
  end process;
end;

library ieee;
use ieee.std_logic_1164.all;
use ieee.std_logic_unsigned.all;
entity mux2 is
  port(inp1,inp2:in std_logic_vector(7 downto 0);
       s:in std_logic;
       outp:out std_logic_vector(7 downto 0));
end entity;
architecture beh of mux2 is
begin
  process(inp1,inp2,s)
  begin
    outp<=(ord('0')&inp1 xor ord('1')&inp2)and s;
  end process;
end;
process(inp1, inp2, s)
begin
  case(s) is
    when '0' =>
      outp<=inp1;
    when '1' =>
      outp<=inp2;
    when others =>
      null;
  end case;
end process;
end;

library ieee;
use ieee.std_logic_1164.all;
entity inv is
  port(a: in std_logic_vector(21 downto 0);
       c: out std_logic_vector(22 downto 0));
end entity;
architecture beh of inv is
  signal d: std_logic_vector(21 downto 0);
begin
  process(a)
  begin
    d<=(not a);
  end process;
  c<=('0'&d);
end;
entity mux4 is
  port(a,b,c,d: in std_logic_vector(22 downto 0);
       s: in std_logic_vector(1 downto 0);
       outp: out std_logic_vector(22 downto 0));
end entity;
architecture beh of mux4 is
begin
  process(a,b,s)
  begin
    case(s) is
      when "00" =>
        outp<=a;
      when "01" =>
        outp<=b;
      when "10" =>
        outp<=c;
      when "11" =>
        outp<=d;
      when others =>
        null;
    end case;
  end process;
end process;
end;

library ieee;
use ieee.std_logic_1164.all;
use ieee.std_logic_unsigned.all;
entity shift2 is
  port(inp: in std_logic_vector(7 downto 0);
       outp: out std_logic_vector(22 downto 0));
end entity;
architecture beh of shift2 is
begin
  process(inp,x1,x2,x3,x4,x5,x6,x7)
  begin
    x1<=(inp & "00");
    x2<=x1+inp;
    x3<=(inp & "000000");
    x4<=x3+inp;
    x5<=(inp & "0000000000000000");
    x6<=(x2 & "0000000000000000");
    x7=x5-x6;
    outp<=x7+x4;
  end process;
end;

library ieee;
use ieee.std_logic_1164.all;
use ieee.std_logic_unsigned.all;
entity swap is
  port(inp1, inp2: in std_logic_vector(22 downto 0);
       outp1, outp2: out std_logic_vector(22 downto 0));
end entity;
architecture beh of swap is
begin
  process(inp1, inp2)
  begin
    outp1<=inp2;
    outp2<=inp1;
  end process;
end;
library ieee;
  use ieee.std_logic_1164.all;
  use ieee.std_logic_unsigned.all;
  use ieee.std_logic_arith.all;
entity cse is
  port(a:in std_logic_vector(7 downto 0);
  c,d:out std_logic_vector(7 downto 0));
end entity;
architecture beh of cse is
  signal s1,s2:std_logic_vector(7 downto 0);
begin
  process(a)
  begin
    s1<="00" & a(7 downto 2);
    s2<="000000" & a(7 downto 6);
  end process;
  c<=a + s1;
  d<=a + s2;
end;
library ieee;
  use ieee.std_logic_1164.all;
entity inv1 is
  port(a:in std_logic_vector(22 downto 0);
  c:out std_logic_vector(22 downto 0));
end entity;
architecture beh of inv1 is
  begin
    process(a)
    begin
      c<='not a';
    end process;
end;
library ieee;
  use ieee.std_logic_1164.all;
entity shift1 is
  port(inp:in std_logic_vector(7 downto 0);
  outp:out std_logic_vector(21 downto 0));
end entity;
architecture beh of shift1 is
  signal x1,x2:std_logic_vector(9 downto 0);
  signal x3,x4:std_logic_vector(13 downto 0);
  signal x5:std_logic_vector(19 downto 0);
  signal x6:std_logic_vector(19 downto 0);
  signal x7:std_logic_vector(21 downto 0);
begin
  process(inp,x1,x2,x3,x4,x5,x6,x7)
  begin
    x1<=(inp & "00");
    x2<x1+inp;
    x3<=(inp & "000000");
    x4<x3+inp;
    x5<=(x2 & "000000000000");
    x6<x2 & "000000000000");
    x7<=(x4 & '0');
    outp=x5+x6+x7;
  end process;
end;
library ieee;
  use ieee.std_logic_1164.all;
entity shift3 is
  port(inp:in std_logic_vector(7 downto 0);
  outp:out std_logic_vector(21 downto 0));
end entity;
architecture beh of shift3 is
  signal x1,x2:std_logic_vector(9 downto 0);
  signal x3,x4:std_logic_vector(13 downto 0);
  signal x5:std_logic_vector(19 downto 0);
  signal x6:std_logic_vector(19 downto 0);
  signal x7:std_logic_vector(21 downto 0);
begin
  process(inp,x1,x2,x3,x4,x5,x6,x7)
  begin
    x1<=(inp & "00");
    x2<x1+inp;
    x3<=(inp & "000000");
    x4<x3+inp;
    x5<="000000000000");
    x6=x5+x2;
    x7<=(x4 & "000000000000");
    outp=x7-x6;
  end process;
end;
library ieee;
  use ieee.std_logic_1164.all;
entity top is
  port(inp:in std_logic_vector(7 downto 0);
  sel1:in std_logic;
  sel2:in std_logic_vector(1 downto 0));
begin
  process(inp,x1,x2,x3,x4,x5,x6,x7)
  begin
    x1<=(inp & "00");
    x2<x1+inp;
    x3<=(inp & "000000");
    x4<x3+inp;
    x5<="000000000000");
    x6=x5+x2;
    x7<=(x4 & "000000000000");
    outp=x7-x6;
  end process;
end;
downto 0);
  outp:out std_logic_vector(22
downto 0));
end entity;
architecture beh of top is
  component cse is
    port(a:in std_logic_vector(7
downto 0);
    c,d:out
std_logic_vector(7 downto 0));
  end component;
  component mux is
    port(a,b:in std_logic_vector(22
downto 0);
    s:in std_logic;
    outp:out std_logic_vector(22
downto 0));
  end component;
  component shift1 is
    port(inp:in std_logic_vector(7
downto 0);
    outp:out std_logic_vector(21
downto 0));
  end component;
  component shift2 is
    port(inp:in std_logic_vector(7
downto 0);
    outp:out std_logic_vector(22
downto 0));
  end component;
  component shift3 is
    port(inp:in std_logic_vector(7
downto 0);
    outp:out std_logic_vector(21
downto 0));
  end component;
  component inv is
    port(a:in std_logic_vector(21
downto 0);
    c:out std_logic_vector(22 downto
0));
  end component;
  component inv1 is
    port(a:in std_logic_vector(22
downto 0);
    c:out std_logic_vector(22 downto
0));
  end component;
  component append is
    port(a:in std_logic_vector(21
downto 0);
    c:out std_logic_vector(22 downto
0));
  end component;
  component swap is
    port(inp1,inp2:in
std_logic_vector(7 downto 0);
    s:in std_logic;
    outp1,outp2:out
std_logic_vector(22 downto 0));
  end component;
  component mux4 is
    port(a,b,c,d:in
std_logic_vector(22 downto 0);
    s:in std_logic_vector(1
downto 0);
    outp:out std_logic_vector(22
downto 0));
  end component;
signal
s1,s2,s3:std_logic_vector(7 downto
0);
signal r1,r2:std_logic_vector(21
downto 0);
signal
t1,inve1,inve2,inve3,inve4:std_logic_vector(22 downto 0);
signal
m1,m2,m3:std_logic_vector(22 downto
0);
begin
  i0:cse port map(inp,s1,s2);
in1:mux2 port
map(s1,s2,sell1,s3);
i2:shift1 port map(s3,r1);
i3:shift2 port map(s3,t1);
i4:shift3 port map(s3,r2);
i5:inv port map(r1,inve1);
i6:inv port map(r2,inve4);
i7:inv1 port map(t1,inve3);
i8:append port map(r1,inve2);
i9:mux port
map(inve1,inve2,sell1,m1);
i10:swap port
map(inve3,inve4,m2,m3);
i11:mux4 port
map(inve1,m1,m2,m3,sell2,outp);
end;
Author Biography

M. Pradeepa was born in Coimbatore in 1987. I complete B.E. Degree in Electronics and communication from Tamilnadu college of engineering, Coimbatore. Currently I am doing M.E. in VLSI Design in SNSCT. I develop this paper as my final year project. I am interested to develop a communication system with the help of VLSI technology. Later I like to expand this paper for complete data transmission modem.